# J-O-Caml (4) <br> jean-jacques.levy@inria.fr pauillac.inria.fr/~levy/qinghua/j-o-caml 

 Qinghua, November 27

## Plan of this class

- polymorphic mutable data
- modules and signatures
- manipulating terms as AST
- reading bitmaps
- ariane 5 story


## Exercices

- Conway sequences - solution 2
\# let conway $\mathrm{x}=$
let rec conway1 $n$ a $x=$ match $x$ with
| [ ] -> if $n>0$ then [1;a] else []
| $b:: y$-> if $a=b$ then conway1 $(n+1) a y$ else
if $n>0$ then $n:: a::$ conway1 1 b y else conway1 1 b y in conway1 $00 \times ;$;


## Polymorphism

- variables and functions are polymorphic in Caml
- type inference gives the principal type (unique most general)
- no need to keep type information at runtime, since strong type checking differs from Lisp, Scheme, Java which keep and check type information
- in Caml, polymorphism only appears around let statements

```
# let id = function x -> x in
    print_int (id 43); print_string (id "jocaml") ;;
43jocaml- : unit = ()
# function x -> x ;;
- : 'a -> 'a = <fun>
```


## Polymorphic mutable data

- mutable values can't have real polymorphic types (see below)
- they are not considered as real values only values have true polymorphic types
- mutable values have "once polymorphic" types

```
# let succ x = x + 1 ;;
val succ : int -> int = <fun>
# let loc = ref (function x -> x) in
    loc := succ; !loc "jocaml" ;;
    Characters 55-63:
        loc := succ; !loc "jocaml" ;;
            ^^^^^^^^
Error: This expression has type string but an expression was expe
cted of type int
# ref (function x -> x) ;;
- : ('_a -> '_a) ref = {contents = <fun>}
```


## Modules and Signatures

- module declaration groups related functions

```
# module FIFO = struct
    type 'a t = {mutable hd: 'a list; mutable tl: 'a list}
    let create() = {hd = [ ]; tl = [ ]}
    let add f a = f.tl <- a :: f.tl
    end;;
module FIFO :
    sig
        type 'a t = { mutable hd : 'a list; mutable tl : 'a list; }
        val create : unit -> 'a t
        val add : 'a t -> 'a -> unit
    end
# |
```

- qualified names to refer to functions and types
\# let f = FIFO.create(); ;
val f : '_a FIFO.t = \{FIFO.hd $=\square$; FIFO.tl $=\square]\}$
\# FIFO.add f 3;
- : unit = ()
\# f;
- : int FIFO.t $=$ \{FIFO.hd $=\square$; FIFO.tl $=[3]\}$


## Modules and Signatures

- implementation may be hidden by forcing signature
\# module FIFO = (struct
type 'a t = \{mutable hd: 'a list; mutable tl: 'a list\}
let create( ) $=\{$ hd $=[\mathrm{]} ; \mathrm{tl}=[\mathrm{]}\}$
let add f $a=$ f.tl <- $a$ :: f.tl
end :
sig
type 'a t
val create : unit -> 'a t
val add : 'a t -> 'a -> unit
end) ;;
module FIFO : sig type 'a t val create : unit -> 'a t val add :
'a t -> 'a -> unit end
- qualified names to refer to functions and types

```
# let f = FIFO.create() ;;
```

val f : '_a FIFO.t = <abstr>
\# FIFO.add f 3; ;

- : unit = ()
\# f ; ;
- : int FIFO.t = <abstr>
\# FIFO.add f "jocaml";
Characters 11-19:
FIFO.add f "jocaml";;


## Modules and Signatures

- modules group set of type, exception, variable, function definitions
- type of a module is its signature
- signature can be restricted by giving it explicitly
- hiding implementation of some types produce abstract types
- several functions may also be hidden (usually auxiliary functions)
- abstract types may have several implementations (FIFO as circular buffers, FIFO as lists)
- if type is abstract, the user of this type will not see differences between implementations
- signatures are described in the Ocaml libraries
- compiled signatures are in files with suffix .cmi
- modules may be nested


## Exercice

- Write remove function which removes the head of the queue in FIFO module (with creation of EmptyQueue exception)
- Give an alternative implementation of FIFOs with circular buffers.
- Give a module definition for addition and multiplication for big numbers (as in exercic lecture 2)


## Reading bitmaps

- function to read bitmaps on standard input +2 useful functions
[format is: width( w ) height( h ) and h lines of w numbers]

```
let ncols = read_int() in
let nlignes = read_int() in
let b = bmap_read nlignes ncols in
bmap_display b;
pause();
let bmap_display b =
    let bi = make_image b in
            draw_image bi margin margin|;;
let pause () =
    match wait_next_event [Button_down] with
        _ -> () ;;
```


## Reading bitmaps

- format is: nlines and ncolumns

```
let bmap_read nlignes ncols =
    let b = Array.make_matrix nlignes ncols 0 in
    for i = 0 to nlignes - 1 do
        let s = read_line() in
        let xs = ref (Str.split (Str.regexp "[\t]+") s) in
        for j = 0 to ncols - 1 do
            let c = int_of_string (List.hd !xs) in
            b.(i).(j) <- rgb c c c;
            xs := List.tl !xs;
        done;
    done;
    b;;
```


## Combien d'objets dans une image? Jean-Jacques Lévy INRIA

## Labeling



16 objects in this picture

## Algorithm

## 1) first pass

- scan pixels left-to-right, top-to-bottom giving a new object id each time a new object is met


## 2) second pass

- generate equivalences between ids due to new adjacent relations met during scan of pixels.


## 3) third pass

- compute the number of equivalence classes


## Complexity:

- scan twice full image (linear cost)
- try to efficiently manage equivalence classes (Union-Find by Tarjan)

